

# Using Parametric Test to determine the Significance of Banded patterns in N-dimensional 0-1 dataset

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# Abstract

**Background**: The identification of banded patterns in N-dimensional (ND) 0-1 dataset is one where all the elements in the dimensions are arranged such that the one's entries are ordered about the center of dimensions. Reordering 0-1 datasets in order to determine banded patterns in data, allows for the identification of interesting pattern that are hidden in the data. The challenge is whether or not the identify banded patterns are significant. Previous work on the significance of banded patterns using parametric test was aimed at 2D and 3D banding algorithms, using a single banding algorithm in each case, which meant no work on the significance of banded patterns using different banding algorithms with different dimensions. **Aim**: To this end, this paper presents a comparison between ND banding algorithms for 2D and 3D. **Method**: The approach is to use parametric test on synthetic data, UCI and the real datasets taken from the cattle tracing system (CTS). The ND banding algorithms considered for 2D are:2D banding, barycentric (BC) and 2D sort, and for 3D are: exact-Euclidean, exact-Manhattan variations and the approximate. **Results**: The experimental results presented shows the significance of banded patterns with p value less than 0.05. However, the post-hoc test result shows a statistically significant difference between 2D banding and BC, 2D banding and 2D sort, exact-Euclidean and exact-Manhattan, exact-Euclidean and the approximate but no significant difference between BC and 2D sort as well as exact-Manhattan and the approximate.

Keywords— Banded pattern, N-dimension, 0-1 data, Parametric test.

# 1. Introduction

The paper presents the techniques for determining the significance of banded patterns in ND 0-1 datasets, that consists of only 0's and 1's. For better understanding, the cell with a one value has a "dot" and zero value empty cell. 0-1 datasets commonly occur in a number of real-world application domains, these are gene and probe mapping in bio-informatics Alizadeh et al., (1995), Cheng and Church (2000), and Myllykangas et al., (2006). Sites and species co-occurrences in paleontology (Atkins et al., (1998) and Mannila et al., (2006)). The advantages of banded patterns include the following; first it gives an insight about the data by clustering records that co-exist. Secondly, it makes visualization of the dataset possible by given a better understanding of the datasets. Thirdly, it compresses the datasets in a way that improves the operation of algorithms that work with 0-1 datasets.

Several studies on Banded pattern identification in ND 0-1 datasets, where (N=2) exists; however, the concept of ND banded patterns as presented in this paper started from the work of Garriga et al., (2008), Garriga et al., (2001) and Makinen and Sirtola (2005) (similarly the work of Abdullahi et al., (2014), Abdullahi et al., (2015), Abdullahi et al., (2018), Abdullahi (2019a), Abdullahi (2018a), Abdullahi (2018b), Abdullahi and Coenen (2018a), Abdullahi and Coenen (2020) in the context of 2D and 3D. The early work on banded patterns identification was based on permutations and was considered NP-complete problem in the work of Papadimitriou (1976). The work presented in this paper is aimed at determining the significance of banded patterns on difference ND banding



algorithms using parametric tests. Given any ND 0-1 dataset with dimensions  $\{\dim_1, \dim_2, \ldots, \dim_n\}$ , we can achieve some banding by rearranging the elements in each dimension. Though in a real-life situation, a perfect banding may not be feasible; but there is a possibility of achieving a best banding. Therefore, the following are the objectives of this paper: (i) to investigate the significance of banded patterns on different ND banding algorithms for (2D and 3D) and (ii) to assess whether the banded patterns obtained from the ND banding algorithms have the same significance or differ in significance. The idea is to find out whether the banded patterns (b) obtained from different ND banding algorithms on datasets (d) after a number of iterations differs significantly or are the same. We let hundred (100) be the expected maximum number of iterations required to achieve banding (This is because the best bandings were all achieved within hundred iterations). We define the null hypothesis (H<sub>0</sub>) by assuming that there is no significant difference between the banding algorithms with a significance level of p = 0.05.

The structure of remaining part of the paper is as follows. The second section presents the materials and methods on different banding algorithms and a brief background review of some relevant work. The third section presents the evaluation results of the ND banding algorithms. The paper discusses the relevant results obtained in section four and finally, section five of the paper concludes with the main findings.

## 2. Materials and Methods

The work presented in this paper is to determine the significance of banded patterns on different banding algorithms that operate in ND 0-1 data sets using parametric test. Banding applications include identifying entities that cooccur in co-occurrence analysis Atkins et al., (1998), Mannila et al., (2006). An example is the paleontological application that combine Neolithic sites with its fossil species. Similarly, Very-Large Scale Integration (VLSI) chip design by Koebe and Knochel (1990), for block alignment problems aimed at using banding concept to identify channels between the circuit component blocks. Finally, in graph drawing, the idea is minimizing the number of edge cross overs, this was achieved using banding concept in Sugiyama et al (1981). The idea of banding emanates from numerical analysis according to Atkinson (2008), where it was used in the context of resolution of linear equations.

The banding concept relates to reorderable matrices, reorderable patterns and bandwidth minimization. Reorderable matrices simply refers to data visualization in 2D. This concept has helped in achieving a better understanding of data as proposed in Bertin (1981), Bertin (1999) and Makinen and Siirtola (2000). In 19th century, an English Egyptologist known as Petrie, used manual reordering approach to understand archaeological datasets as mentioned in Liiv (2010). For reorderable pattern, the idea is to arrange a 2D matrix column wise and row wise so that hidden pattern of interests can be identified (Aykanat et al., (2004) and Deutsch and Martin (1971)).

The approach finds a pattern P, revealed by arranging a given 2D matrix column wise and row wise, this however are not necessarily 0-1 datasets as such the pattern P may be of any form. The problem with this concept is how to rearrange 2D datasets to find the pattern P. Similarly, minimization of Bandwidth according to Cheng (1973a), Cheng (1973b) is the process of minimizing the bandwidth of 2D sparse matrix by arranging it row wise and column wise such that the one's entries are as close as possible along the main diagonal. However, it is worth noting that there are several studies on methods of data reordering in different applications domain. Such example is the Barycentric (BC) algorithm proposed by Makinen and Sirtola (2005), which is directed at 2D data, originally the concept support the drawing of graphs.

The algorithm finds permutation of 2D matrix row wise and column wise such that the one's entries converge about the diagonal using a barycentric measure. The BC algorithm finds the average position of each one's entries column wise/row wise. In this paper, the BC algorithm was considered one of the 2D algorithms used to determine the significance of banded patterns. Similarly, the 2D sort method, is an approach used to reorder 2D



matrix in order to detect patterns along the main diagonal of the matrix Bertin (1999), Harri and Eekki (2005). The 2D sort method works by computing the sum of each given 2D matrix column wise and row wise and then sorted the matrix accordingly. For the 3D banding algorithms, the approximate algorithm finds an approximate banding using dimension pairing; while the exact algorithm finds an exact banding for the entire data space.

The two metrics considered for exact banding algorithms are the Euclidean and Manhattan distance proposed by Abdullahi et al., (2016). There is limited work on the significance of banded patterns, for instance the work of Abdullahi and Coenen (2018b), Abdullahi and Hassan (2020), proposed the statistical significance testing for banded patterns using Gaussian distribution. The idea was to apply Gaussian distribution on randomly generated 2D datasets without banding and then used it confirmed whether the obtained banded patterns are significant or not with respect to their distance from mean.

The experimental results indicate the statistical significance of banding with at least one to two standard deviations from the mean. Additionally, the findings also establish the possibility of generating Gaussian distribution curves with ranges of one density values. Although, the results were promising however, the limitation was that it requires generating the Gaussian distribution curve for each of the data sets. Abdullahi (2019b) also proposed the use of student t-statistic and chi-square test to determine the statistical significance of 3D banded patterns. The experimental results presented shows 5% level of significant, and the paper only consider the exact-Euclidean metric of the 3D banding algorithm. Similarly, Abdullahi et al (2021), proposed the used of three parametric tests: student t-test, paired student t-test and ANOVA test to determine the significance of 2D banding Algorithms. The results from the experimental analysis shows a significant difference between the banding Algorithms.

The identification of banded patterns in ND 0-1 datasets requires arranging all elements in the dimensions to identify a best banding (or an approximate if no best exist). The idea presented by Abdullahi et al., (2015), use the "banding scores" (BS) concept. BS is a score value obtained by summing the one's entries associated with each index in the dimension and divide by the maximum number of the one's entries. The datasets comprise of set of dimensions (DIM) that consist of n number of dimensions, where  $DIM = \{\dim_1, \dim_2, \ldots, \dim_n\}$ . Note that the dimension sizes are not equal. The dimension i represents a sequence of k index values  $\{e_{i1}, e_{i2}, \ldots, e_{ik}\}$ , i the dimension identifier and  $e_{ij}$  represents item e in dimension i with respect to index j. Every location has 0's and 1's. The distribution of the data entries is based on the application domain. Every one entry represents a set of coordinates (c1, c2, ..., cn), where n is the number of dimensions and c1  $\in \dim_1, c2 \in \dim_2$  and so on.

For 2D,  $DIM = \{\dim_i, \dim_j\}$ , meaning computing the banding scores for dimension i with respect to dimension j and vice versa, by simply adding the index values for y(x) and dividing by the maximum indexes. For 2D datasets, the banding score ( $bs_{iq}$ ,) is computed as shown in Eq. 1.

$$bs_{iq} = \frac{\sum_{i=1}^{i=|M|} m_i \in M \times n_i \in N}{\sum_{i=1}^{i=|M|} (k_p - i + 1) \times n_i \in N'}$$
(1)

Where:  $M = \{m_1, m_2, \dots, m_k\}$  list of dimension j indexes for the one's entry that features index q in dimension i,  $N = \{n_1, n_2, \dots, n_k\}$  list of the number of one's entries at each location corresponding to list l, N' represents list similar to N but rearranged in descending order to the number of one's entry at each location and  $k_j$  is the size of dimension j. Converting Eq. 1 to address ND data, Eq. 2 is used to calculate the banding scores for ND.



$$bs_{iq} = \frac{\sum_{i=1}^{i=|L|} dis \tan ce(l_i \in L) \times n_i \in N}{\sum_{i=1}^{i=|MAXDIST|} (\max Dist_p \in MAXDIST) \times n_i \in N'}$$
(2)

Where: *L* represents list of locations in the data space of size n less than 1, and n = |DIM|, *MAXDIST* represents list of maximum distances in descending order where |MAXDIST| is equivalent to |L|, *N* and *N'* as previously defined. The function distance () returns a distance to the origin of the datasets with one less than the original data space presented as coordinates set  $(s_1, s_2, s_3 \cdots, s_{n-1})$ .

For the 3D approximate banding algorithm, the banding score for index k in dimension i with respect to dimension j, denoted as  $bs_{ijk}$ , is calculated using Eq. 3.

$$bs_{ijk} = \frac{\sum_{k=1}^{k=|N|} n_k}{\sum_{l=1}^{l=|N|} (|dim_l| - p + 1)}$$
(3)

Where N is the set of dimension j indexes representing one's entry whose dimension i coordinate is equivalent to k  $(N = \{n_1, n_2, ...\})$ . The Overall Score (OS) for the whole configuration is computed using Eq. 4, where  $OS_i$  is the overall scores for dimension *i*.

$$OS = \frac{\sum_{i=1}^{i=|\text{dim}|} OS_i}{|\text{dim}|} \tag{4}$$

The  $OS_i$  value is calculated using Eq. 5. Here the banding is weighed with respect to its index location, as the aim is to arrange the one's entries from the top corner of the data space.

$$OS_{i} = \frac{\sum_{j=1}^{j=|p_{i}|} bs_{ij} \times (p_{i} - k + 1)}{\frac{p_{i}(p_{i} + 1)}{2}}$$
(5)

#### 2.1 The ND Banding Algorithm

The pseudocode for the ND banding algorithm is presented in Listing 1. The inputs in (lines 1 and 2) are: 0-1 data set D, consisting set of tuples of the form (c1, c2, ...) indicating the location of each one's entry in the datasets; and  $DIM = \{\dim_1, \dim_2, \dots, \dim_n\}$  set of dimensions in D. The output is an arranged data set D with minimum OS value. The algorithm iterates over the datasets. For every iteration, the algorithm rearranges the indexes in DIM, using the scoring concept in Eq. 1, 2 and 3 to identify a revised DIM (Step 3). This revised DIM is then used to arrange D to give D' (Step 9). A new value of OS is computed (using Eq. 3 and 4). If the new OS value (OS<sub>new</sub>) is less than the current OS values. The algorithm terminates with the previous data configuration and the value OS. Hence, D, DIM and OS are updated (Steps 9 to 11) and the algorithm continues.

INPUT: 0-1 dataset D with DIM, where DIM = {dim<sub>1</sub>, dim<sub>2</sub>, ..., dim<sub>n</sub>} OUTPUT: N rearrange to minimize OS Begin: Step 1: OS = 1.0 Step 2: LOOP

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**Step 3**: DIM' = The rearranged DIM using Eq. 1, 2 and 3 **Step 4**: D' = The rearranged D according to DIM' **Step 5**:  $OS_{new}$  = calculated OS values for D using Eq. 4 and 5 **Step 6**: IF ( $OS_{new} \ge OS$ ) THEN **Step 7**: break **Step 8**: ELSE **Step 9**: D = D' **Step 10**: DIM = DIM' **Step 11**: OS =  $OS_{new}$  **Step 12**: END IF **Step 13**: END LOOP **Step 14**: Terminate with D and OS **End** 

Listing 1. Generic ND banding Algorithm

The approximate variation, of the ND banding algorithm pairs the dimensions to compute the banding scores as presented in Eq. 3. We replace Steps 3 and 4 in Listing 1, with the pseudocode in Listing 2. The Approximate ND algorithm operates by considering all pairings ij in the dimension. For each pairing, the banding score  $b_{Sijk}$  for e index k in dimension i is computed with respect to dimension j (Step 4), and the indexes in dimension i are arranged according to the computed banding score values (Step 6) and then D (Step 8).

**Step 1:** FOR i = 1 to i = DIM| - 1 DO **Step 2:** FOR j = i + 1 to j = |DIM| and  $j \neq i$  DO **Step 3:** FOR k = 1 to  $p = |p_i|$  DO **Step 4:**  $b_{s_{ijk}} =$  calculate the banding score for index k in dimension i with respect to dimension j using Eq. 3 **Step 5:** END FOR **Step 6:** DIM = rearrange dimension i according to their  $b_{s_{ijk}}$  values **Step 7:** D = rearrange D according to DIM<sub>i</sub> **Step 8:** END FOR **Step 9:** END FOR

Listing 2. The Approximate ND banding Algorithm

For the exact ND variation, we replace Steps 3 and 4 in Listing 1, with pseudocode given in Listing 3. The exact ND banding algorithm iterates over the entire datasets computing the banding scores for index q in dimension i using Eq. 2. For each dimension, it rearranges the indexes in the dimension according to the  $bs_{ij}$  values (Step 5), then reconfigure D (Step 6).

**Step 1:** FOR i = 1 to i = |DIM| DO **Step 2:** FOR j = 1 to j = ki DO **Step 3**:  $bs_{ij} =$  calculate banding score for index q in dimension i using Eq. 2 **Step 4:** END FOR



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Step 5:  $DIM = rearrange DIM_i$  according to  $bs_{ij}$  values Step 6: D = rearrange D according to  $DIM_i$ Step 7: END FOR

Listing 3. The exact ND banding Algorithm

The 2D banding algorithm iterate over the 2D datasets by calculating the banding scores for index q in dimension i using Eq. 1 (Steps 3 to 5) as presented in Listing 4. The calculated banding scores are used to rearrange the indexes in the dimension i (Step 6). The banding scores for each index q in dimension j is then calculated using Eq. 1 (Steps 7 to 9). The calculated banding scores are used to rearrange the indexes in the dimension j (Step 10). The new OS values is calculated using Eq. 4 and 5 (Step 11), if the new OS value is greater than the previous OS value (Step 12), then algorithm exits and return M and OS (Steps 15 and 16). Otherwise, the algorithm repeats and exits if no further change.

INPUT: n x m (0-1) matrix M OUTPUT: Reordered matrix row wise and column wise **BEGIN:** Step 1: OS = 1.0Step 2: LOOP Step 3: FOR ALL  $j \in \{1 \cdots \dim_i\}$  DO Step 4:  $Bs_{ij}$  = calculate banding score for dimension i using Eq. 1 Step 5: END FOR Step 6: M = Rearranged dimension i according to the banding score Step 7: FOR ALL  $i \in \{1 \cdots \dim_i\}$  DO Step 8:  $bs_{ii}$  = calculate banding score for dimension j using Eq. 1 Step 9: END FOR Step 10: M' = Rearranged dimension j according to the banding score Step 11:  $OS_{new}$  = calculated OS value for M using Eq. 4 and 5 Step 12: IF  $(OS_{new} \ge OS)$  THEN Step 13: break Step 14: ELSE Step 15: M = M'Step 16:  $OS = OS_{new}$ Step 17: END IF Step 18: END LOOP Step 19: EXIT with M and OS **END** 

Listing 4. The 2D Banding Algorithm

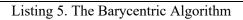


The BC algorithm first calculates the BC for each row *i*, and then arranged the rows in ascending order of BC values. The matrix then transposes and recomputes the BC values for the columns and arrange it in ascending order of values. The process repeats until no more changes. This BC compute the Average Moment (AM) as shown in Eq. 6. Where  $b_{ij}$  is the *j*<sup>th</sup> entry for the row /column *i*, and *n* number of columns / rows.

$$AM = \frac{\sum_{j=1}^{n} jb_{ij}}{\sum_{i=1}^{n} b_{ij}}$$
(6)

Listing 5 presents the pseudocode for the BC algorithm. The input is a 0-1 data set M. The output is a reordered matrix M. The BC computes the AM for all rows in M (Steps 2 to 4), then reorders the rows in ascending order of the AM value (Step 5). The algorithm then transposes the matrix  $M^{T}$  (Step 6) and iterates until no more change (Step 7). It exits with M (Step 10).

INPUT: An n x m (0-1) matrix M OUTPUT: Reordered matrix row wise and column wise BEGIN: Step 1: LOOP Step 2: FOR each row i  $\epsilon$  M DO Step 3: Calculate AM values for row i using Eq. 6 Step 4: END FOR Step 5: M' = rearrange M with row wise according to AM values in ascending order Step 6: M<sup>T</sup> = M' transpose Step 7: Repeat process on M' until no more change Step 8: M' = M Step 9: END LOOP Step 10: EXIT with M END



The 2D sort algorithm iteratively reorder a 2D matrix so as to identify banded patterns, the one's entries known as "black areas" in Bertin (1981) are arranged about the main diagonal. The 2D sort algorithm works by calculating the sum of rows and columns and then arranges them row wise / column wise in ascending order of the row /column sum. The 2D sort algorithm then calculate the row / column sum as shown in Eq. 7. Where N<sub>ij</sub> is the position j of the cell for row/column *i*, and *n* is the number of columns /rows.

$$S = \sum_{j=1}^{j=n} N_{ij} \tag{7}$$

The pseudocode for the 2D sort algorithm is presented in Listing 6. The input is a 0-1 data set M. The output is a rearranged matrix M. The 2D sort algorithm calculates the sum for all rows in M (Steps 2 to 4), then reorders them row wise in ascending order of their sum (Step 5). The algorithm then calculates the sum for all columns in M' (Steps 6 to 8), then reorders them in column wise in ascending order of their sum (Step 9). The process is repeated until no more changes (Step 10). It then terminates with M (Step 12).

INPUT: An n x m (0-1) matrix M OUTPUT: Rearranged matrix M row wise and column wise BEGIN: Step 1: LOOP



Step 2: FOR EACH row i  $\epsilon$  M DO Step 3: Calculate row sum for row i using Eq. 7 Step 4: END FOR Step 5: M' = rearrange M row wise in ascending order of the row sum Step 6: FOR EACH column j  $\epsilon$  M DO Step 7: Calculate column sum for column j using Eq. 7 Step 8: END FOR Step 9: M'' = rearrange M column wise in ascending order of the column sum Step 10: Repeat process on M'' until no more change Step 11: END LOOP Step 12: EXIT with M END

Listing 6. The 2D sort Algorithm

### 3. Results

This section presents the experimental results obtained from the ND banding algorithms. We conduct the experiments using: (i) 2D randomly generated datasets in Coenen (2003), (ii) 2D datasets taken from University of California Irvine data repository Blake and Merz (1998), (iii) 3D randomly generated datasets in Coenen (2003) and (iv) 3D datasets taken from the Cattle Tracing System repository in Nohuddin et al., (2010). The evaluation objectives are to compare the statistically significant difference between the 2D banding algorithms: (i) 2D banding, (ii) BC, (iii) 2D sort and the 3D banding algorithms: (i) exact-Euclidean, (ii) exact-Manhattan and (iii) the approximate. Finally, to determine whether these algorithms are similar in significance and which differs. The parametric test considered are: (i) student t-test, (ii) Paired student t-test and (iii) ANOVA Parametric test is a statistical test that assumes parameters with known distributions about population. The parametric test was considered in the paper because it uses mean value to measure the central tendency, which relies on statistical distributions of data, and also has more chances of accuracy.

After the significant difference between the algorithms are tested, the post-hoc test is performed to determine which of these algorithms differs from the other. In this paper, Benjamin-Hochberg post hoc test proposed by Gupta in (2013) was adopted and used to determine the algorithms that differs in significance. The post-hoc test method works by initially testing the hypothesis under consideration and compute the test-statistic, which is then use to obtain a p-value from the distribution. Table 1 presents the statistics test and their respective p-values in bold fonts for (i) 2D banding with BC, (ii) 2D banding with 2D sort, (iii) BC with 2D sort, as well as (iv) 2D banding, BC with 2D sort. Similarly, Table 2 presents the statistic test and the p-values in bold fonts for (i) exact-Euclidean with the approximate, (iii) exact-Euclidean with the approximate as well as (iv) exact-Euclidean, exact-Manhattan with the approximate. Table 3 then presents the post-hoc test results for p = 0.05 in bold fonts respectively. Figures 1 and 2 is the bar chart representation of 2D random datasets with the number of iterations. Similarly Figures 3 and 4 is the bar chart representation of the 3D random datasets with the number of iterations and the CTS datasets with number of iterations respectively.

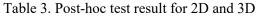
Parametric test	2D Banding algorithms	Random datasets		UCI datasets	
		Statisti c test	p-value <0.05	Statisti c test	p-value < 0.05
Student t-test	2D banding with BC	2.4056	0.0271	2.4839	0.0230
	BC with 2-D Sort	3.1716	0.0052	4.4385	0.0003
	2D banding with 2D sort	2.4251	0.0260	2.5310	0.0209
table-value	C C	1.8330	0.05	1.8330	0.05
Paired Student t-	2D banding with BC	3.0937	0.0128	2.3551	0.0429
test	BC with 2D Sort	3.3797	0.0081	5.0825	0.0006
	2D banding with 2D sort	3.2299	0.0103	2.7664	0.0218
	C C	2.2620	0.05	2.2620	0.05
table-value					
ANOVA	2D banding, BC with 2D sort	7.7799	0.0021	11.336	0.0002
table-value		4.4100	0.05	4.4100	0.05

Table 2. 3D Parametric Statistic test result

Parametric test	<b>3D</b> banding algorithms	3D Random datasets		CTS datasets	
		Statisti c test	p- value <0.05	Statisti c test	p-value < 0.05
Student t-test	exact-Euclidean with exact-Manhattan	2.9858	0.0079	2.6060	0.0178
	exact-Euclidean with the Approximate	6.8963	1.89 <sup>e-0.6</sup>	5.3033	4.83 <sup>e-0.5</sup>
	exact-Manhattan with the Approximate	2.7021	0.0145	2.9395	0.0087
table-value		1.8330	0.05	1.8330	0.05
Paired Student	exact-Euclidean with exact-Manhattan	2.4046	0.0395	2.7688	0.0217
t-test	exact-Euclidean with the Approximate	7,0601	5.91 <sup>e-05</sup>	4.3528	0.0018
	exact-Manhattan with the Approximate	2.8900	0.0178	2.7618	0.0220
table-value		2.2620	0.05	2.2620	0.05
ANOVA	Exact-Euclidean, exact-Manhattan with the Approximate	18.239	9.72 <sup>e-06</sup>	14.373	5.616 <sup>e-05</sup>
table-value		4.4100	0.05	4.4100	0.05



2D banding algorithms	P=0.05	Reject (H <sub>0</sub> )	3D banding algorithms	P=0.05	Reject (H <sub>0</sub> )
2D banding with BC	-0.9157	True	exact-Euclidean with exact- Manhattan	-0.4437	True
2D banding with 2D sort	-2.8157	True	exact-Euclidean with the approximate	-2.1437	True
BC with 2D sort	0.0843	False	exact-Manhattan with the approximate	0.1563	False



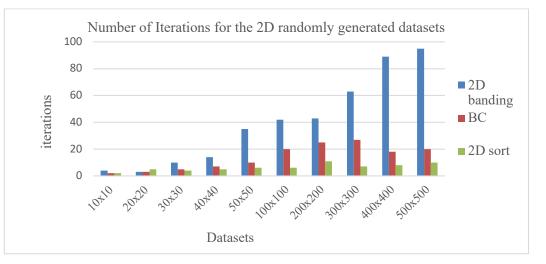


Figure 1. 2D random data with Number of iterations

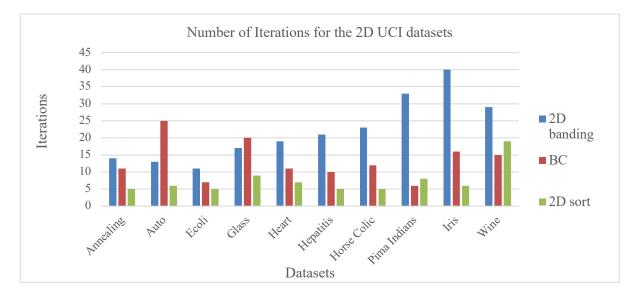


Figure 2. 2D UCI datasets with Number of iterations



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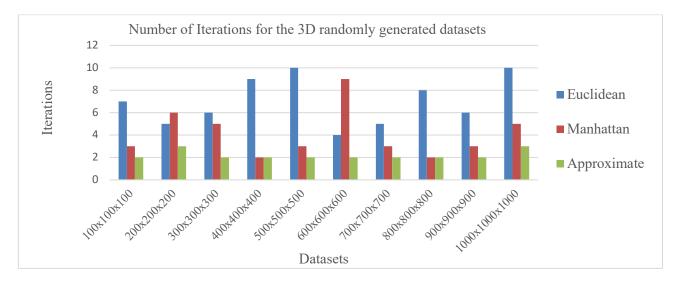


Figure 3. 3D random datasets with Number of iterations

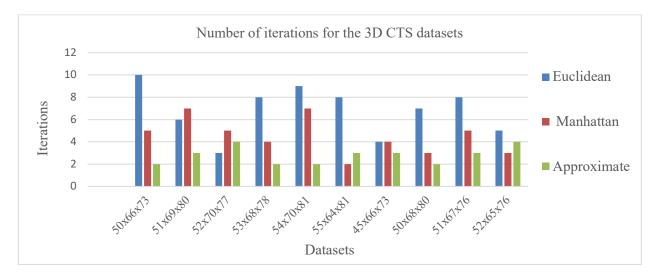


Figure 4. 3D CTS datasets with Number of iterations.

# 4. Discussion

The experimental results presented in Tables 1 and 2 above shows the significance of banded patterns with respect to 0.05 p value for the 2D and 3D respectively. From table 1, the result for the 2D random datasets and the UCI datasets are presented. The p value results for the 2D banding with BC, 2D banding with 2D sort and BC with 2D sort are: **0.0271**, **0.0052**, **0.026**, **0.0230**, **0.0003** and **0.0209** less than 0.05 using the student t-test we thereby reject  $H_0$  that state there is no significant difference between the banding algorithms and accept  $H_A$  that assume significant difference between the 2D random datasets and the UCI datasets using paired



student t-test, the p value results were 0.0128, 0.0081, 0.0103, 0.0429, 0.0006 and 0.0218 less than 0.05, as such we reject  $H_0$  and accept  $H_A$ . Furthermore, the results for ANOVA test are 0.0021 and 0.0002 respectively still less than 0.05, thus  $H_0$  is rejected and  $H_A$  accepted. Table 2 presents the p values results for 3D random and CTS datasets: the student t -test results are 0.0079, 1.89e-0.5, 0.0145, 0.0178, 4.83e-0.5, 0.0087 which are less than 0.05, hence  $H_0$  is rejected,  $H_A$  accepted. The Paired student t-test results are: 0.0395, 5.91e-0.5, 0.0178, 0.0217, 0.0018 and 0.0220 still less than 0.05, thus we reject  $H_0$  and accept  $H_A$ . Similarly, the ANOVA test results are 9.72e-0.6 and 5.616e-0.5 less than 0.05,  $H_0$  rejected but  $H_A$  accepted. Table 3 present the post-hoc test, from the table, the experimental results for 2D banding with BC is -0.9157 less than 0.05  $H_0$  rejected is False. The results for exact-Euclidean with exact-Manhattan is -0.4437 less than 0.05  $H_0$  rejected is True, exact-Euclidean with the approximate is -2.1437 less than 0.05  $H_0$  rejected is True and finally, exact-Manhattan with the approximate is 0.1563 greater than 0.05  $H_0$  rejected is False.

## 5. Conclusion

The ND banding algorithms have been presented. Specifically, the 2D banding algorithms are: 2D banding, BC and 2D sort and the 3D banding algorithms are: exact-Euclidean, exact-Manhattan and the approximate respectively. The experimental results presented confirmed the following findings: (i) All the banding algorithms tested were statistically significant with p values less than 0.05 for randomly generated, UCI and the CTS datasets respectively for the three parametric test considered, (ii) The post-hoc test result shows significant difference between: (i) 2D banding with BC, (ii) 2D banding with 2D sort, (iii) exact-Euclidean with exact-Manhattan, (d) exact-Euclidean with the approximate with their p values less than 0.05. However, it fails to show a significant difference between (i) BC with 2D sort and (ii) exact-Manhattan with the approximate as their p values were greater than 0.05 respectively. For future work, the authors intend to investigate visualization of the ND banded patterns.

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